#### The Complexity of the Game of



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Each card has 4 attributes:

- Symbol
- Number
- Color
- Shading



- Symbol
  - Oval
  - Diamond
  - Squiggle



- Number
  - One
  - Two
  - Three



- Color
  - Red
  - Purple
  - Green



- Shading
  - Solid
  - Stripped
  - Blank



There are  $3^4 = 81$  different cards in total (one for each combination of values).



- Deal 12 cards;
- Find a *valid set:* set of 3 cards such that, for each attribute, the values are either *all the same* or *all different*.

Valid set: 3 cards with values for each attribute being either all the same or all different.



- All have same color;
- all have same symbol;
- all have same number;
- all have different shadings.

Valid set: 3 cards with values for each attribute being either all the same or all different.



- All have different colors;
- all have different symbols;
- all have different numbers;
- Il have same shading.

Valid set: 3 cards with values for each attribute being either all the same or all different.



- All have different colors;
- all have different symbols;
- all have different number;
- all have different shadings.

This is not a valid set!



- All have same colors;
- all have different symbols;
- × only 2/3 have same number;
- all have different shadings.

### Interesting fact

 Each pair of cards has a unique 3<sup>rd</sup> with which they constitute a valid set.



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### Naive way to find a set

• For each pair of cards, check whether the card that completes the valid set exists among the rest.



#### Generalization: k-1SET

- Input: *m* cards, *n* attributes, *k* values
- Question: Does there exist a valid set of k cards with all values the same or all values different?
- In the original game m=12, n=4 and k=3.

## Complexity Results for k-1SET

- For *k* unbounded:
  - $n = 2 \rightarrow P^1$
  - *n* ≥ 3 → NP-Complete<sup>1</sup>
- For *n* unbounded:
  - $k = 2 \rightarrow \text{trivial}$
  - k parameter  $\rightarrow XP^{1}$

#### k-1SET parameterized by k is W-hard (reduction applies to a version of *n-Dim Matching*)

1. Chaudhuri et al: On the complexity of the game of set (manuscript 2003).

• Naive algorithm works for multi-round version, seeking the max number of co-existing valid sets.



• However, finding the max number of valid sets with card removal is not so easy (even for *k*=3).



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For k = 3 and *unbounded n*:

- Find a single valid set is in P.
- Find max # *r* of co-existing valid sets is in P.
- max-3-rSET: Find the max # r of existing valid sets with card removal.
- min-3-rSET: Find the min # r of valid sets that destroy all other valid sets.

Construct a 3-uniform hypergraph: vertices  $\leftrightarrow$  cards, hyperedges  $\leftrightarrow$  valid sets.

- max-3-*r*SET: restriction of 3-Set Packing to SET hypergraphs).
- min-3-*r*SET: restriction of Independent Edge Dominating Set to SET hypergraphs).



Both restrictions remain NP-hard. min-3-rSET parameterized by r is FPT

## Interesting fact for 3-SET

Assign each value a number among {0,1,2}. Then, each card is an *n*-tuple of ternary coordinates.

In a valid set, if we add together corresponding coordinates we always get 0 (mod 3).

#### <u>Proof</u>

For a particular attribute:

- all values the same:  $3*i = 0 \pmod{3}$ ;
- all values different:  $0+1+2 = 0 \pmod{3}$ .

#### Conclusions

- We studied the complexity of single- and multi-round variations of the game of SET.
- We established connections of the game with Multi-Dimensional Matching, Set Packing, Independent Edge Dominating Set.

# TII NK YOU!

Michael Lampis, Valia Mitsou: The Computational Complexity of the Game of Set and its Theoretical Applications. *Latin 2014.*